Property Estimation ++

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Including past work with J. Acharya, H. Das, A.T. Suresh, K. Viswanathan

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Outline

Property estimation

Prior work

Plug-in estimators

Maximum likelihood

Profile maximum likelihood

Simple unified approach

Proof elements

Results

Discrete Distributions

Discrete support set \mathcal{X}

 ${\text{heads, tails}} = {\text{h, t}} \qquad \mathbb{Z}$

Distribution p over \mathcal{X} , probability p_x for $x \in \mathcal{X}$

$$\begin{aligned} p_x &\geq 0 \qquad \sum_{x \in \mathcal{X}} p_x = 1 \\ p &= (p_{\mathsf{h}}, p_{\mathsf{t}}) \qquad p_h = .6, \ p_t = .4 \end{aligned}$$

 ${\cal P}$ collection of distributions

 $\mathcal{P}_{\mathcal{X}}$ all distributions over \mathcal{X}

 $\mathcal{P}_{\{h, t\}} = \{(p_h, p_t)\} = \{(.6, .4), (.4, .6), (.5, .5), (0, 1), \ldots\}$

Distributions Property

 $f:\mathcal{P}_{\mathcal{X}}\to\mathbb{R}$

Maps distribution to real value, also called functional

Shannon entropy	H(p)	$\sum_x p_x \log \frac{1}{p_x}$		
Support size	S(p)	$\sum_x \mathbb{1}_{p_x > 0}$		
Support coverage	$S_m(p)$	$\sum_x (1 - (1 - p_x)^m)$		
Expected $\#$ distinct symbols in m samples				
Distance to uniformity	$L_{uni}(p)$	$\sum_{x} \left p_x - \frac{1}{ \mathcal{X} } \right $		
Rényi entropy	$H_{\alpha}(p)$	$\frac{1}{1-\alpha}\log\left(\sum_x p_x^\alpha\right)$		
Highest probability	$\max(p)$	$\max\left\{p_x: x \in \mathcal{X}\right\}$		

Many applications

Symmetric Properties

f invariant under label permutations $H(p) \quad H_{\alpha}(p) \quad S(p) \quad S_m(p) \quad L_{uni}(p) \quad \max(p)$ Non-symmetric: f depends on labels

 $p_{\mathsf{h}} = \frac{p_{\mathsf{h}}}{p_{\mathsf{t}}} = p_{\mathsf{h}} \cdot p_{\mathsf{t}}, \text{ if } |\mathcal{X}| > 2$

Additive Properties

$$f(p) = \sum_{x} f(p_{x})$$

$$S(p) \coloneqq \sum_{x} \mathbb{1}_{p_{x}>0}$$

$$H(p) = \sum_{x} p_{x} \log \frac{1}{p_{x}}$$

$$S_{r}(p)$$

$$L_{uni}(p)$$

Non-additive

 $H_{\alpha}(p) \coloneqq \frac{1}{1-\alpha} \log\left(\sum_{x} p_{x}^{\alpha}\right)$ $\max(p) \coloneqq \max\left\{p_{x} : x \in \mathcal{X}\right\}$

Most results apply to additive symmetric properties

Property Estimation

Given: support set \mathcal{X} , property fUnknown: $p \in \mathcal{P}_{\mathcal{X}}$ Estimate: f(p)Entropy of English words Given: $\mathcal{X} = \{\text{English words}\}, f = H$, unknown: p, estimate: H(p)# species in habitat Given: $\mathcal{X} = \{\text{bird species}\}, f = S$, unknown: p, estimate: S(p)

Learn from examples

Observe n independent samples X^n = $X_1,\ldots,X_n \sim p$ Estimate f(p)



Estimator: $f^{\text{est}} : \mathcal{X}^n \to \mathbb{R}$ Estimate: $f^{\text{est}}(X^n)$

Plug-in Estimators

Simple two-step estimators Use X^n to derive estimate $p^{\text{est}}(X^n)$ of pPlug-in $f(p^{\text{est}}(X^n))$ to estimate f(p)If as $n \to \infty$, $p^{\text{est}}(X^n) \to p$, then $f(p^{\text{est}}(X^n)) \to f(p)$

What is the simplest p^{est} ?

Empirical Estimator

$n \, \operatorname{samples}$

 $N_x \ \# \ times \ x \ appears$

 $p_x^{\mathsf{emp}} \coloneqq \tfrac{N_x}{n}$

Entropy estimation

$$\begin{split} \mathcal{X} &= \{a, b, c\} \qquad p = (p_a, p_b, p_c) = (.5, .3, .2) \\ \text{Estimate } H(p) \text{ from } n = 10 \text{ samples} \\ X^{10} &= c, a, b, a, b, a, b, a, b, c \\ p^{\mathsf{emp}} &= (.4, .4, .2) \\ H^{\mathsf{emp}}(X^{10}) &= H(.4, .4, .2) \end{split}$$

Best-known, most widely-used distribution estimator

Relatively easy to analyze

Best Estimator?

Min-max formulation

Given: Property f, collection \mathcal{P} of distributions over \mathcal{X} *n* i.i.d. samples X^n from unknown $p \in \mathcal{P}$ Property value f(p) – unknown Estimator's value $f^{\text{est}}(X^n)$ Estimator's absolute loss $|f^{\mathsf{est}}(X^n) - f(p)|$ Expected loss $L_f(f^{\text{est}}, p, n) := \mathbb{E}_{X^n \sim p} |f^{\text{est}}(X^n) - f(p)|$ Worst-case loss $L_f(f^{\text{est}}, \mathcal{P}, n) \coloneqq \max_{p \in \mathcal{P}_{\mathcal{V}}} L_f(f^{\text{est}}, p, n)$ Minimum worst-case loss $L_f(\mathcal{P}, n) \coloneqq \min_{f^{est}} L_f(f^{est}, \mathcal{P}, n)$

Abbreviation

Symmetric properties

 $\mathcal{P}_{\mathcal{X}}$ all distributions over \mathcal{X}

Dependence on \mathcal{X} only through $k = |\mathcal{X}|$

H over {cat, dog} same as over {ma, shu}

 $L_f(\mathcal{P}_{\mathcal{X}}, n) \rightarrow L_f(k, n)$

Prior work: Min-max Error up to Constant Factors

Property	Base function	$L(f^{emp},k,n)$	L(k,n)	
Entropy ¹	$p(x)\log \frac{1}{p(x)}$	$\frac{k}{n} + \frac{\log n}{\sqrt{n}}$	$\frac{k}{n\log n} + \frac{\log n}{\sqrt{n}}$	
Supp. coverage ²	$(1 - (1 - p(x))^r)$	$r\exp\left(-\Theta\left(\frac{n}{r}\right)\right)$	$r \exp\left(-\Theta\left(\frac{n\log n}{r}\right)\right)$	
Power sum ^{3 4}	$p(x)^{lpha}$, $lpha \in (0, rac{1}{2}]$	$\frac{k}{n^{lpha}}$	$\frac{k}{(n\log n)^{\alpha}}$	
	$p(x)^{\alpha}$, $\alpha \in (\frac{1}{2}, 1)$	$\frac{k}{n^{\alpha}} + \frac{k^{1-\alpha}}{\sqrt{n}}$	$\frac{k}{(n\log n)^{\alpha}} + \frac{k^{1-\alpha}}{\sqrt{n}}$	
Dist. to uniform ⁵	$ p(x) - rac{1}{k} $	$\sqrt{\frac{k}{n}}$	$\sqrt{\frac{k}{n \log n}}$	
Support size ⁶	$\mathbb{1}_{p(x)>0}$	$k \exp\left(-\Theta\left(\frac{n}{k}\right)\right)$	$k \exp\left(-\Theta\left(\sqrt{\frac{n\log n}{k}}\right)\right)$	

References: P03, VV11a/b, WY14/19, JVHW14, AOST14, OSW16, ADOS17, JVW18

n to $n \log n$ when comparing the worst-case performances

$$\begin{array}{l} {}^{1}n\gtrsim k \text{ for empirical; } n\gtrsim k/\log k \text{ for minimax} \\ {}^{2}n\gtrsim r \text{ for empirical; } n\gtrsim r/\log r \text{ for minimax} \\ {}^{3}\alpha\in(0,\frac{1}{2}]:\ n\gtrsim k^{1/\alpha} \text{ for empirical; } n\gtrsim \frac{k^{1/\alpha}}{\log k} \text{ and } \log k\gtrsim \log n \text{ for minimax} \\ {}^{4}\alpha\in(\frac{1}{2},1):\ n\gtrsim k^{1/\alpha} \text{ for empirical; } n\gtrsim \frac{k^{1/\alpha}}{\log k} \text{ for minimax} \\ {}^{5}n\gtrsim k \text{ for empirical; } n\gtrsim k/\log k \text{ and } \log k\gtrsim \log n \text{ for minimax} \\ {}^{6}\text{consider } \mathcal{P}_{\ge 1/k} \text{ instead of } \mathcal{P}_{\mathcal{X}};\ k\log k\gtrsim n\gtrsim k/\log k \text{ for minimax} \end{array}$$

Why is Empirical Suboptimal?

Intuitive, simple

Why does it work at all?

Maximum Likelihood

For i.i.d. $p \in \mathcal{P}_{\mathcal{X}}$, the probability of observing $x^n \in \mathcal{X}^n$

$$p(x^n) \coloneqq \Pr_{X^n \sim p}(X^n = x^n) = \prod_{i=1}^n p(x_i)$$

Maximum likelihood estimator: $x^n \rightarrow \text{dist.} p \text{ maximizing } p(x^n)$

$$p^{\mathsf{ml}}(x^n) = \arg \max_p p(x^n)$$
$$p^{\mathsf{ml}}(h,t,h) = \arg \max_{p_h} p_h^2 \cdot (1-p_h) \quad \rightarrow \quad p_h = 2/3, \ p_t = 1/3$$
Identical to empirical estimator – always

Good: distribution that best explains observation

Sub-optimal for all properties in table

ML / EF work well for small alphabets large sample

Overfit data when alphabet is large relative to sample size

Can we improve?

iid: Do not care about order

Symmetric properties: Do not care about specific values (h,h,t), (t,t,h), (h,t,h), (t,h,t), (t,h,h), (h,t,t) same entropy Care only: # of elements appearing any given number of times Three samples: 1 element appeared once, 1 element appeared twice Profile: $\varphi = \{1,2\}$

Profile maximum likelihood (PML)

Profile $\varphi(x^n)$ of x^n is the multiset of its symbol frequencies

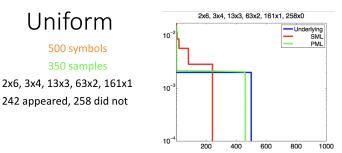
 $\begin{array}{l} x^n = a \, b \, a \, c \, c \, d \, e \implies a \, c \text{ appears twice, } b \, d \, e \text{ appear once} \\ \implies \varphi(x^n) = \{2,2,1,1,1\} \end{array}$

Probability of observing a profile φ when sampling from p is

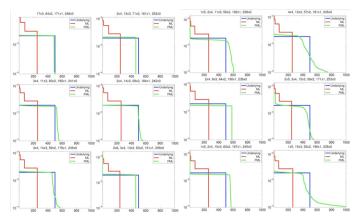
$$p(\varphi) \coloneqq \sum_{y^n:\varphi(y^n)=\varphi} p(y^n) = \sum_{y^n:\varphi(y^n)=\varphi} \prod_{i=1}^n p(y_i)$$

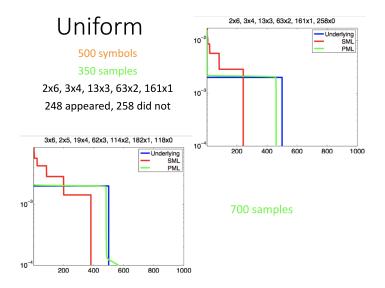
[OSVZ04] Profile maximum likelihood maps x^n to

$$p_{\varphi(x^n)}^{\mathsf{ml}} \coloneqq \underset{p \in \mathcal{P}_{\mathcal{X}}}{\operatorname{argmax}} p(\varphi(x^n))$$

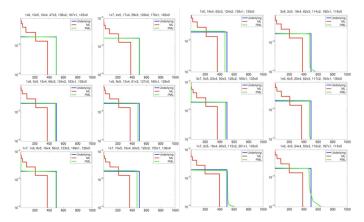


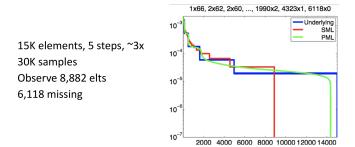
U[500], 350x, 12 experiments

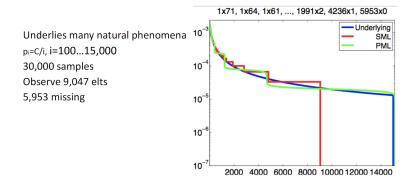




U[500], 700x, 12 experiments







1990 Census - Last names

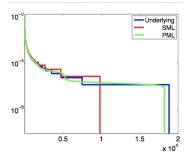
SMITH	1.006	1.006	1
JOHNSON	0.810	1.816	2
WILLIAMS	0.699	2.515	3
JONES	0.621	3.136	4
BROWN	0.621	3.757	5
DAVIS	0.480	4.237	6
MILLER	0.424	4.660	7
WILSON	0.339	5.000	8
MOORE	0.312	5.312	9
TAYLOR	0.311	5.623	10
AMEND	0.001	77.478	18835
ALPHIN	0.001	77.478	18836
ALLBRIGHT	0.001	77.479	18837
AIKIN	0.001	77.479	18838
ACRES	0.001	77.480	18839
ZUPAN	0.000	77.480	18840
ZUCHOWSKI	0.000	77.481	18841
ZEOLLA	0.000	77.481	18842
	0.000	//.101	10012

18,839 names 77.48% population ~230 million

1990 Census - Last names

18,839 last names based on ~230 million

35,000 samples, observed 9,813 names

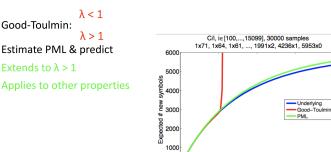


Coverage (# new symbols)

Zipf distribution over 15K elements

Sample 30K times

Estimate: # new symbols in sample of size λ * 30K



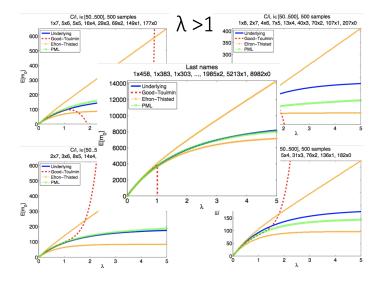
0

1

2

3

4



Proof Elements

Is unlikely likely?

Upper bound probability of observing unlikely outcomes

```
p{:}\ {\rm distribution}\ {\rm over}\ {\mathcal Z}
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 $\delta > 0$

 $z \in \mathcal{Z}$ is δ -unlikely if $p(z) \leq \delta$

$$\begin{split} &\Pr(\text{observing a } \delta - \text{ unlikely outcome}) = \sum_{z \in \mathcal{Z}_{\leq \delta}} p(z) \leq \sum_{z \in \mathcal{Z}_{\leq \delta}} \delta = \\ &\delta \cdot |\mathcal{Z}_{\leq \delta}|. \end{split}$$

Competitiveness of PML

Consider the problem of symmetric property estimation Φ_n : collection of profiles associated with samples of size nLemma Suppose $\hat{f}: \Phi_n \to \mathbb{R}$ is such that for all $p \in \mathcal{P}_X$,

$$\Pr_{\varphi \sim p}(|\hat{f}(\varphi) - f(p)| > \varepsilon) < \delta,$$

then the PML plug-in estimator satisfies [ADOS17]

$$\Pr_{\varphi \sim p} \left(|\boldsymbol{f}(\boldsymbol{p}_{\varphi}^{\mathsf{ml}}) - f(p)| > 2 \cdot \varepsilon \right) < \delta \cdot \exp(3\sqrt{n})$$

Competitiveness of PML

Proof: Consider any $p \in \mathcal{P}_{\mathcal{X}}$ $\Phi_{\geq \delta}^{n} := \{\varphi \in \Phi_{n} : p(\varphi) \geq \delta\}$ For $\varphi \in \Phi_{\geq \delta}^{n}$: $|\hat{f}(\varphi) - f(p)| \leq \varepsilon$ (condition in the lemma) $p_{\varphi}^{\mathsf{ml}}(\varphi) \geq p(\varphi) \geq \delta$, hence $|\hat{f}(\varphi) - f(p_{\varphi}^{\mathsf{ml}})| \leq \varepsilon$ Triangle inequality: $|f(p_{\varphi}^{\mathsf{ml}}) - f(p)| \leq 2\varepsilon$

Therefore,

$$\Pr_{\varphi \sim p} \left(|\boldsymbol{f}(\boldsymbol{p}_{\varphi}^{\mathsf{ml}}) - f(p)| > 2\varepsilon \right) \leq \Pr_{\varphi \sim p} (\varphi \notin \Phi_{\geq \delta}^{n}) \leq \delta \cdot |\Phi_{n}|$$

Finally, $|\Phi_n|$ is exactly the number of partitions of integer n, which $\leq \exp(3\sqrt{n})$ by the well-known result* of Hardy and Ramanujan

*Hardy, G. H. and Ramanujan, S. "Asymptotic Formulae in Combinatory Analysis." Proc. London Math. Soc. 17, 75-115, 1918.

Sample Complexity Formulation

```
p an unknown distribution in \mathcal{P}_{\mathcal{X}}
Given an i.i.d. sample X^n \sim p
Estimate f(p) by estimator \hat{f}
Min-max sample complexity n_f(|\mathcal{X}|, \varepsilon, \delta)
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minimum n necessary to
ensure |\hat{f}(X^n) - f(p)| \le \varepsilon with probability \ge 1 - \delta
for every p \in \mathcal{P}_{\mathcal{X}}
```

Equivalent to result in table

The Broad Optimality of PML [HO19a]

Profile maximum likelihood (PML) is a unified timeand sample-optimal approach to four fundamental problems: additive property estimation, Rényi entropy estimation, uniformity testing, and sorted distribution estimation.

Hao, Y., & Orlitsky, A. (2019). The Broad Optimality of Profile Maximum Likelihood.

Additive property estimation

Theorem For every f in a broad class of symmetric additive properties, including all Lipschitz properties, any \mathcal{X} , $p \in \mathcal{P}_{\mathcal{X}}$, and $n \ge n_f(|\mathcal{X}|, \varepsilon, 1/3)$, if $\varepsilon \ge n^{-0.1}$,

$$\Pr\left(\left|f\left(p_{\varphi(X^{4n})}^{\mathsf{ml}}\right) - f(p)\right| > 5\varepsilon\right) \le \exp(-\sqrt{n}).$$

Can use APML [CSS19], approximating PML in near linear time.

Observations

Prior work either:

Used different estimators for different properties

Applied a plug-in estimator for only few properties

(A)PML apply to all additive Lipschitz properties and more

Essentially strengthens original table

Runs in near-linear time

Additional results

α -Rényi entropy estimation

For integer $\alpha > 1$, PML plug-in has optimal $k^{1-1/\alpha}$ sample complexity

For non-integer $\alpha > 3/4$, (A)PML plug-in improves best-known results

Sorted distribution estimation

Under ℓ_1 distance, (A)PML yields optimal $\Theta(k/(\varepsilon^2 \log k))$ sample complexity for sorted distribution estimation

Uniformity testing: $p = p_u$ v.s. $|p - p_u| \ge \varepsilon$; complexity $\Theta(\sqrt{k}/\varepsilon^2)$

Tester below is sample-optimal up to logarithmic factors of \boldsymbol{k}

Input: parameters k, ε , and a sample $X^n \sim p$ with profile φ If any symbol appears $\ge 3 \max\{1, n/k\} \log k$ times, return 1 If $||p_{\varphi}^{\mathsf{ml}} - p_u||_2 \ge 3\varepsilon/(4\sqrt{k})$, return 1; else, return 0

Thank You!